# Social Network Analysis Lecture 2-Introduction Graph Theory 

Donglei Du<br>(ddu@unb.ca)

Faculty of Business Administration, University of New Brunswick, NB Canada Fredericton E3B 9Y2

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## What is Graph theory?

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- A graph may be undirected, meaning that there is no distinction between the two vertices associated with each edge, or directed, meaning that its arcs may be directed from one vertex to another.


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- A graph may be undirected, meaning that there is no distinction between the two vertices associated with each edge, or directed, meaning that its arcs may be directed from one vertex to another.
- This is an extremely brief introduction of graph theory.


## The Seven Bridges of Königsberg:



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- Q: Is there a Eulerian trail through the city that would cross each bridge once and only once? (The second (undirected) graph represents the bridge network!)
- A: an (connected) undirected graph has an Eulerian trail if and only if at most two vertices have odd degree.


## The Seven Bridges of Königsberg:



- Q: Is there a Eulerian trail through the city that would cross each bridge once and only once? (The second (undirected) graph represents the bridge network!)
- A: an (connected) undirected graph has an Eulerian trail if and only if at most two vertices have odd degree.
- Leonhard Euler (15 April 1707-18 September 1783, Swiss Mathematician) in 1735 laid the foundations of graph theory and prefigured the idea of topology by studying this problem.


## The Erdös coauthor network and Erdös number:



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- Paul Erdös (26 March 191320 September 1996, Hungarian mathematician): one of the most prolific publishers of papers in mathematical history, comparable only with Leonhard Euler; Erdös published more papers, mostly in collaboration with other mathematicians, while Euler published more pages, mostly by himself.


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- Find your own Erdös Number


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http://www.ams.org/mathscinet/collaborationDistance.html
- If you want to know more about Erdös Number,try here: http://www.oakland.edu/enp/compute/


## Six Degrees of Kevin Bacon:



## Six Degrees of Kevin Bacon:



- Kevin Norwood Bacon (July 8, 1958-) is an American actor and musician


## Types of graphs

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- Undirected vs directed:



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- Undirected vs directed:

- Undirected networks: coauthorship network, actor network, Königsberg Bridges Network, Facebook friendship network


## Types of graphs

- Undirected vs directed:

- Undirected networks: coauthorship network, actor network, Königsberg Bridges Network, Facebook friendship network
- Directed networks: URLs on the www, phone calls, Retweet network


## Types of graphs

## Types of graphs

- Simple vs multigraph: lops or multiedges



## Types of graphs

## Types of graphs

- Unweighted vs weighted:



## Types of graphs

- Unweighted vs weighted:

- Mobile phone calls, Collaboration network


## Important graphs

## Important graphs

- Regular graph


## Important graphs

- Regular graph
- Complete graph


## Important graphs

- Regular graph
- Complete graph
- Path


## Important graphs

- Regular graph
- Complete graph
- Path
- Cycle


## Important graphs

- Regular graph
- Complete graph
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## Important graphs

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- Euler graph


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- ...


## Layout

## Representation of graphs using different data

 structure (Hi, Computer Science guys)
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- Edge list
- Adjacency list
- Laplace matrix


## Adjacency matrix: undirected graph



## Adjacency matrix: undirected graph

- Adjacency matrix: $A_{i j}=1$ iff there is a link between $i$ and $j$.

$$
\left[\begin{array}{lllllll} 
& 1 & 2 & 3 & 4 & 5 & 6 \\
1 & 0 & 1 & 0 & 0 & 1 & 0 \\
2 & 1 & 0 & 1 & 0 & 1 & 0 \\
3 & 0 & 1 & 0 & 1 & 0 & 0 \\
4 & 0 & 0 & 1 & 0 & 1 & 1 \\
5 & 1 & 1 & 0 & 1 & 0 & 0 \\
6 & 0 & 0 & 0 & 1 & 0 & 0
\end{array}\right]
$$

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\end{array}\right]
$$

- The adjacency is symmetric for undirected graph.


## R code based on package igraph: generate graph from adjacency matrix

```
library(igraph)
#Generate graph object from adjacency matrix
adjm_u<-matrix(
    c(0, 1, 0, 0, 1, 0,
1, 0, 1, 0, 1, 0,
0, 1, 0, 1, 0, 0,
0, 0, 1, 0, 1, 1,
1, 1, 0, 1, 0, 0,
0, 0, 0, 1, 0, 0), # the data elements
    nrow=6, # number of rows
    ncol=6, # number of columns
    byrow = TRUE) # fill matrix by rows
g_adj_u <- graph.adjacency(adjm_u, mode="undirected")
tkplot(g_adj_u)
```


## Adjacency matrix: directed graph



## Adjacency matrix: directed graph

- Adjacency matrix: $A_{i j}=1$ iff there is a link from $j$ to $i$

$$
\left[\begin{array}{llllll} 
& 1 & 2 & 3 & 4 & 5 \\
1 & 0 & 0 & 0 & 0 & 0 \\
2 & 1 & 0 & 0 & 1 & 0 \\
3 & 0 & 1 & 0 & 1 & 0 \\
4 & 0 & 1 & 0 & 0 & 1 \\
5 & 0 & 1 & 0 & 0 & 0
\end{array}\right]
$$

- Note the direction of the edge runs from the second index to the first: counter-intuitive, but convenient mathematically!


## R code based on package igraph: generate graph from adjacency matrix

```
library(igraph)
#Generate graph object from adjacency matrix
adjm_d<-matrix(
    c(0, 1, 0, 0, 0,
0, 0, 1, 1, 1,
0, 0, 0, 0, 0,
0, 1, 1, 0, 0,
0, 0, 0, 1, 0), # the data elements
    nrow=5, # number of rows
    ncol=5, # number of columns
    byrow = TRUE) # fill matrix by rows
g_adj_d <- graph.adjacency(adjm_d, mode="directed")
tkplot(g_adj_d)
```


## Edge list: undirected graph



## Edge list: undirected graph

- Edge list:
$(1,2)$
$(1,5)$
$(2,3)$
$(2,5)$
$(3,4)$
$(4,5)$
$(4,6)$


## R code based on package igraph: generate undirected graph from edge list

```
library(igraph)
#Generate graph object from edge list
el_u <- matrix( c(1, 2, 1, 5, 2,3,2,5,3,4,4,5,4,6), nc=2,
g_el_u <- graph.edgelist(el_u,directed=FALSE)
tkplot(g_el_u)
```


## Edge list: directed graph



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## R code based on package igraph: generate directed graph from edge list

```
library(igraph)
#Generate graph object from edge list
el_d <- matrix( c(1, 2, 2, 3, 2,4,4,2,2,5,4,3,5,4), nc=2,
g_el_d <- graph.edgelist(el_d,directed=TRUE)
tkplot(g_el_d)
```


## Adjacency list: undirected graph



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- Adjacency list:


| node | neighbor |
| :--- | :--- |
| $1:$ | 2a |
| $2:$ | 13 |

## Adjacency list: undirected graph

- Adjacency list:

| node | neighbor |
| :--- | :--- |
| $1:$ | 25 |
| 2: | 1335 |
| 3: | 244 |
| 4: | 356 |
| 5: | 124 |
| 6: | 4 |

- Easier to work with if network is large and parse, and quick in retrieving all neighbors for a node


## Adjacency list: directed graph



## Adjacency list: directed graph



- Adjacency list:

| node | inneighbor | outneighbor |
| :--- | :--- | :--- |
| $1:$ |  | 2 |
| 2: | 1,4 | $3,4,5$ |
| $3:$ | 2,4 |  |
| 4: | 2,5 | 2,3 |
| $5:$ | 2 | 4 |

## Layout

## Graph-theoretic concepts

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- Density


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- Degree, indegree and outdegree


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- Density
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- Path and cycles
- Distance, diameter
- Components
- Clustering coefficient


## Degree for undirected graph



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- Degree of any node $i$ : the number of nodes adjacent to $i$. It can be calculated from the three different representations discussed earlier.


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- Degree of any node $i$ : the number of nodes adjacent to $i$. It can be calculated from the three different representations discussed earlier.
- Every loop adds two degrees to a node.

| node | degree |
| :--- | :--- |
| 1 | 2 |
| 2 | 3 |
| 3 | 2 |
| 4 | 3 |
| 5 | 3 |
| 6 | 1 |

## R code based on package igraph: degree

```
degree(g_adj_u)
```


## Indegree and outdegree for directed graph



## Indegree and outdegree for directed graph



- Indegree of any node $i$ : the number of nodes destined to $i$.


## Indegree and outdegree for directed graph



- Indegree of any node $i$ : the number of nodes destined to $i$.
- Outdegree of any node $i$ : the number of nodes originated at $i$.


## Indegree and outdegree for directed graph



- Indegree of any node $i$ : the number of nodes destined to $i$.
- Outdegree of any node $i$ : the number of nodes originated at $i$.
- Every loop adds one degree to each of the indegree and outdegree of a node.

| node | indegree | outdegree |
| :--- | :--- | :--- |
| 1 | 0 | 1 |
| 2 | 2 | 3 |
| 3 | 2 | 0 |
| 4 | 2 | 2 |
| 5 | 1 | 1 |

## R code based on package igraph: degree

```
degree(g_adj_d, mode="in")
degree(g_adj_d, mode="out")
```


## Walk and Path



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- Walk: a walk is a sequence of nodes in
 which each node is adjacent to the next one


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- In a directed network, the walk can follow only the direction of an arrow.
- A path is a walk without passing through the same link more than once (e.g. 1, 2, 5, 4, 3.).
- Cycle: A path with the same start and end node (e.g., 1, 2, 5)


## Distance



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- The Distance between every pair of nodes are

|  | 1 | 2 | 3 | 4 | 5 | 6 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 1 | 0 | 1 | 2 | 2 | 1 | 3 |
| 2 | 1 | 0 | 1 | 2 | 1 | 3 |
| 3 | 2 | 1 | 0 | 1 | 2 | 2 |
| 4 | 2 | 2 | 1 | 0 | 1 | 1 |
| 5 | 1 | 1 | 2 | 1 | 0 | 2 |
| 6 | 3 | 3 | 2 | 1 | 2 | 0 |

## Distance

- The distance $d_{i j}$ (shortest path, geodesic path) between two nodes $i$ and $j$ is the number of edges along the shortest path connecting them.
- The Distance between every pair of nodes are

|  | 1 | 2 | 3 | 4 | 5 | 6 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 1 | 0 | 1 | 2 | 2 | 1 | 3 |
| 2 | 1 | 0 | 1 | 2 | 1 | 3 |
| 3 | 2 | 1 | 0 | 1 | 2 | 2 |
| 4 | 2 | 2 | 1 | 0 | 1 | 1 |
| 5 | 1 | 1 | 2 | 1 | 0 | 2 |
| 6 | 3 | 3 | 2 | 1 | 2 | 0 |

- For directed graphs, the distance from one node $A$ to another $B$ is generally different from that from $B$ to $A$.


## R code based on package igraph: distance

shortest.paths(g_adj_u)

## Number of walks of length $k$

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## Number of walks of length $k$

- If $A$ is the adjacency matrix of the directed or undirected graph $G$, then the matrix $A^{k}$ (i.e., the matrix product of $k$ copies of
$A)$ has an interesting interpretation:
- the entry in row $i$ and column $j$ gives the number of (directed or undirected) walks of length $k$ from vertex $i$ to vertex $j$.
- This implies, for example, that the number of triangles in an undirected graph $G$ is exactly the trace of $A^{3} / 3$ !.


## Connectivity and connected components for undirected graph



Figure: A graph with three connected components

## Connectivity and connected components for undirected graph

- A graph is connected if there is a path between every pair of nodes.

Figure: A graph with three connected components

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- A graph is connected if there is a path between every pair of nodes.
- A connected component is a subgraph in which any two nodes are connected to each other by paths, and which is connected to no additional nodes in the original graph.


## Connectivity and connected components for undirected graph



Figure: A graph with three connected components

- A graph is connected if there is a path between every pair of nodes.
- A connected component is a subgraph in which any two nodes are connected to each other by paths, and which is connected to no additional nodes in the original graph.
- Largest Component: Giant Component


## Connectivity and connected components for undirected graph



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- A connected component is a subgraph in which any two nodes are connected to each other by paths, and which is connected to no additional nodes in the original graph.
- Largest Component: Giant Component
- Bridge: an edge whose deletion increases the number of connected components.


## Connectivity and connected components for directed graph



Figure: A graph with three strongly connected components

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- A directed graph is weakly connected
if it is connected by ignoring the edge directions


## Connectivity and connected components for directed graph

- A directed graph is strongly connected if there is a path from any node to each other node, and vice versa.
- A directed graph is weakly connected if it is connected by ignoring the edge directions
- A strongly connected component of a directed graph $G$ is a subgraph that is strongly connected, and is maximal with this property: no additional edges or vertices from $G$ can be included in the subgraph without breaking its property of being strongly directed.


## Clustering coefficient



## Clustering coefficient

- For any given node $A$ and two randomly selected nodes $B$ and $C$ :

$$
\begin{aligned}
C C(A) & =\mathbb{P}(B \in N(C) \mid B, C \in N(A)) \\
& =\mathbb{P}(\text { two randomly selected friends of } A \text { are friends }) \\
& =\mathbb{P}(\text { fraction of pairs of } A \text { 's friends that are linked to each other }) .
\end{aligned}
$$

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& =\mathbb{P}(\text { fraction of pairs of } A \text { 's friends that are linked to each other }) .
\end{aligned}
$$

- For example, in the Figure above,

| node | CC |
| :---: | :---: |
| A | 1 |
| B | $1 / 3$ |
| C | $1 / 3$ |
| D | 0 |
| E | 0 |

## R code based on package igraph: Clustering coefficient

\# First generate the graph from edge list el_cc <- matrix ( c("A", "B", "A","C", "B", "C", "B", "E", "D", "E","C","D"), nc=2, byrow=TRUE)
g_el_cc <- graph.edgelist(el_cc, directed=FALSE)
\# Then calculate CC
transitivity(g_el_cc, type="localundirected")

## Duncan Watts - The Myth of Common Sense

## Duncan Watts - The Myth of Common Sense



- http://www. youtube.com/watch?feature=player_ detailpage\&v=D9XF0QOzWM0

References I

